
Mathematical Methods for Computer Science I

Fall 2016

Series 12 – Hand in before Monday, 19.12.2016 - 13.00

Reminder: Wednesday, 21.12.2016: exercises session; Thursday, 22.12.2016: lecture

1. Give context-free grammars that generate the following languages.
 - a) The language of admissible sequences of *if* and *else*. An if-clause can appear alone with no else-clause, or it may be balanced by exactly one else-clause.
 - b) $\{w \in \{a, b, c\}^* \mid w \text{ contains at least three } b\text{'s}\}$.
Hint: think first about the minimal number of productions needed.
 - c) The set $\{a^i b^j c^k \in \{a, b, c\}^* \mid i \neq j \text{ or } j \neq k\}$, that is, the set of strings of *a*'s followed by *b*'s followed by *c*'s, such that there is either a different number of *a*'s and *b*'s or a different number of *b*'s and *c*'s, or both.
Hint: for an unequal number of *a*'s and *b*'s use productions that either repetitively generate an equal number of *a*'s and *b*'s or produce any number of *a*'s or any number of *b*'s and then move on to another production.
2. Let G_1 and G_2 be context-free grammars. Prove that $L(G_1) \cup L(G_2)$ is a context-free language.
3. Design a PDA which accept the following language of all strings of 0's and 1's with an unequal number of 0's and 1's.
4. Consider the language $L = \{0^k 1^k 2^k \mid k \geq 0\}$. Use the pumping lemma for context-free languages to show that L is not context-free.